

DTU



Algorithmic Algebraic Geometry

Project

Find an efficient algorithm to speed up real root counting for univariate tetranomials with high probability. Approach will be by approximating A-discriminant contours in a new way.

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- **Degenerate Roots:** Degenerate roots help describe transitions in number of real roots and closeness to degeneracy governs hardness of numerical solving.
- **Topological Behavior:** More generally, degenerate roots describe transitions in the isotopy type of a (varying) real algebraic surface.

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- **Discretizing Partial Differential Equations:** In certain physical modelling problems, one is trying to approximate the solutions of a very complicated differential equation. So one then uses a numerical scheme to approximate the solution, and this usually involves expanding into a basis of polynomials. Getting information about the solution a PDE can then be reduced to solving a structured polynomial system, many times, with varying coefficients, over the real numbers.

Easy Example

Using the following tetranomial:

$$c_0 + c_1x + c_2x^2 + c_3x^3 \quad (1)$$

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yields a manageable discriminant:

$$-27c_0^2c_3^2 + 18c_0c_1c_2c_3 - 4c_0c_2^3 - 4c_1^3c_3 + c_1^2c_2^2$$

Harder Example

Using a nastier tetranomial:

$$c_0 + c_1x^3 + c_2x^5 + c_3x^{19} \quad (2)$$

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yields a nastier result!:

$$\begin{aligned}
 & 1978419655660313589123979 c_0^{16} c_3^5 + 6093825838807983035604992 c_0^{12} \\
 & c_1^3 c_2^2 c_3^4 - 416630859061143640782400 c_0^{10} c_1 c_2^7 c_3^3 + 4136784303514917397331968 c_0^8 c_1^6 c_2^4 c_3^3 - \\
 & 168062625401816003641344 c_0^6 c_1^{11} c_2 c_3^3 + 546553895696624329228288 c_0^6 c_1^4 c_2^9 c_3^2 + \\
 & 304059692558924048760832 c_0^4 c_1^9 c_2^6 c_3^2 + 9103573347707241984000 c_0^4 c_1^2 c_2^{14} c_3 + \\
 & 24410972524327076888576 c_0^2 c_1^{14} c_2^3 c_3^2 - 1103132840914428362752 c_0^2 c_1^7 c_2^{11} c_3 + \\
 & 34725021329868800000 c_0^2 c_2^{19} + 498062089990157893632 c_1^{19} c_3^2 - \\
 & 48896735641570639872 c_1^{12} c_2^8 c_3 + 1200096737160265728 c_1^5 c_2^{16}
 \end{aligned}$$

Moving Forward...

We need a better way to plot the zero sets of complicated polynomials!
We will use the clever Horn-Kapranov Uniformization to reduce the dimension of the parameter space!

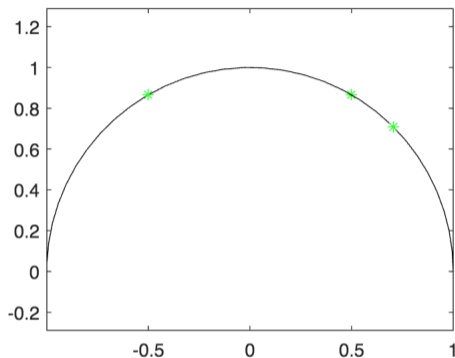
Horn-Kapranov Uniformization

A way to efficiently parameterize discriminant varieties. For $A = [a_1 \ a_2 \ a_3 \ a_4]$, let \hat{A} be the 2×4 matrix defined by appending a row of 1s to the top of A and let B in $\mathbb{Z}^{4 \times 2}$ be any matrix whose columns form a basis for the right nullspace of \hat{A} . Then the (logarithmic, reduced) Horn-Kapranov Uniformization for A is the function

$$\xi_A([\lambda_1 : \lambda_2]) := (\text{Log} | [\lambda_1, \lambda_2] B^T |) B$$

which defines a map from $\mathbb{P}_{\mathbb{R}}^1$ to \mathbb{R}^2 .

Horn-Kapranov Uniformization II



For nicer plots, we use: $(\lambda_1 : \lambda_2) = (\cos\theta, \sin\theta)$

This brings our plots from: $[\lambda_1 : \lambda_2] \in P_R^1$

to: $(\lambda_1 : \lambda_2) \in \text{Unit Semi-Circle}$.

Amoeba

If f is any polynomial in $\mathbb{C}[x_1, \dots, x_n]$ then its amoeba is the set

$$\{(\log |x_1|, \dots, \log |x_n|) \mid f(x_1, \dots, x_n) = 0, x_i \in \mathbb{C} \setminus \{0\}\}$$

.

Amoeba



Figure: This is the Amoeba for $1 + x_1 + x_2$.

Image obtained from: [https://en.wikipedia.org/wiki/Amoeba_\(mathematics\)](https://en.wikipedia.org/wiki/Amoeba_(mathematics))

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- The boundary of the last amoeba is defined by the graphs of “simple” transcendental function, e.g., $y = \text{Log}(1 + e^x)$.
- Deciding if a rational point lies on or near such a curve gets us into interesting problems involving Diophantine approximation!

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- Alternative: Approximate each amoeba boundary curve by a piecewise linear curve.
- Such curves can be extracted from the Horn-Kapranov Uniformization.
- Do they work well with random polynomials/points?
- Experiments show: So-so...

Experimentation!!!

The ultimate goal of our experimentation is to understand how well tropical discriminant chambers approximate true sign chambers.



Matlab Code Round 1

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- Tests a set number of i.i.d. random points to see if they are within that chamber
- Yields an accuracy percentage

Using a Tropical, Linear Approximation:

We use the piecewise function $y=0$ and $y=x$ to approximate the curve of the amoeba.

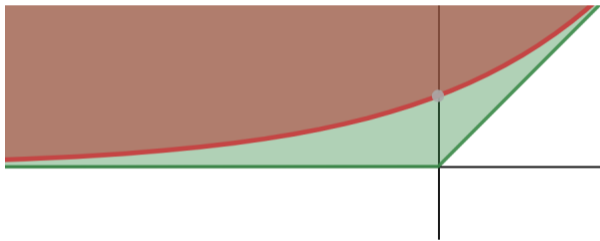


Figure: $y \geq \log(1 + e^x)$ and $y \geq 0$ and $y \geq x$

Results are so-so:

Testing 1000 i.i.d. random points:

Trials:	1	2	3	4	5
%:	62%	65%	63%	60%	65%

Testing 10,000 i.i.d. random points:

Trials:	1	2	3	4	5
%:	63%	63%	65%	64%	64%

Testing 100,000 i.i.d. random points:

Trials:	1	2	3	4	5
%:	64%	63%	64%	64%	64%

Complexity Issue for Chamber Membership

- Deciding a polynomial inequality, involving a polynomial of degree d in n variables with coefficients all of absolute value $\leq H$, at an input rational point $p = (a_1/b_1, \dots, a_n/b_n)$, is a highly non-trivial problem!

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- Deciding a polynomial inequality, involving a polynomial of degree d in n variables with coefficients all of absolute value $\leq H$, at an input rational point $p = (a_1/b_1, \dots, a_n/b_n)$, is a highly non-trivial problem!
- We will use a little trick to get around this! We will change x and y to logarithmic values to yield a more manageable equation to test our inequalities.

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- The arcs of the discriminant contour are defined by linear combinations of logarithms.
- Approximate each arc by just 2 logarithms: This should also yield easier Diophantine approximation.



Matlab Code Round 2

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- Code runs through two loops to apply the Horn-Kapranov Uniformization
- The associated amoeba to the polynomial is given along with each quadrant

Example:

We look at a family of polynomials

Canonical slice of $\text{Nabla}_A(\mathbb{R})$, plotted on log paper, for the family
 $c_1 + c_2x^7 + c_3x^{22} + c_4x^{55}$

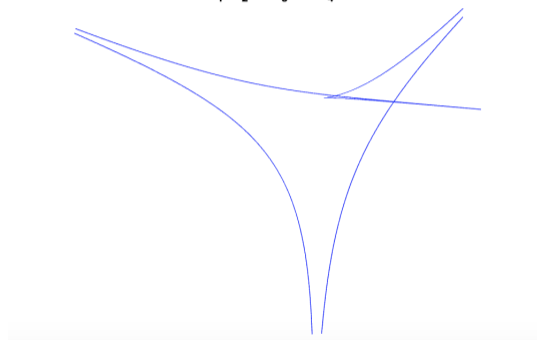
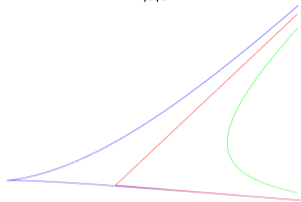


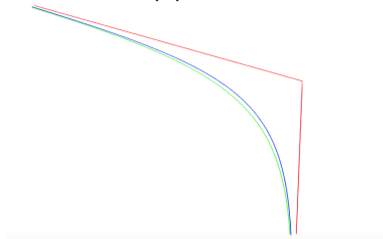
Figure: This is the Ameoba for $c_1 + c_2x^7 + c_3x^{22} + c_4x^{55}$.

Much Closer Approximations (for the most part):

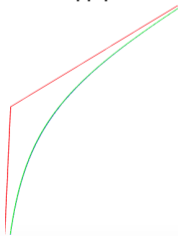
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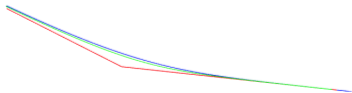
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Testing the Coefficient Space

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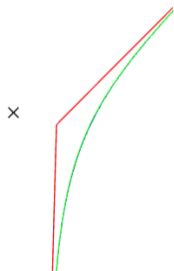
- Setting the coefficients of the polynomial plots a point in the quadrant!

Testing the Coefficient Space

- Setting the coefficients of the polynomial plots a point in the quadrant!
- This shows us where the polynomial lies in coefficient space!

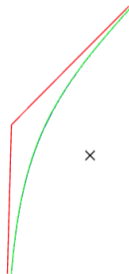
Plotting Polynomials as Points:

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$$-1 - x^7 + x^{22} - x^{55}$$

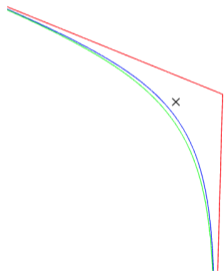
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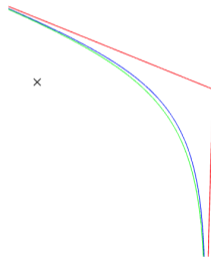
$$-2 - 2x^7 + 10x^{22} - x^{55}$$

More Points:

$$\begin{array}{c} + + + \\ + - + \end{array}$$

$$\begin{array}{c} + + + \\ + - + \end{array}$$


$$-1 + x^7 + x^{22} - x^{55}$$



$$-20 + 10x^7 + 2x^{22} - 10x^{55}$$

Obtaining a Real Root Count:

Plugging the previous polynomial examples into Maple will give the real roots.

Results:

$$f1 := -1 - x^7 + x^{22} - x^{55};$$

`realroot(f1);`

$$f2 := -2 - 2x^7 + 10x^{22} - x^{55};$$

`realroot(f2);`

$$f3 := -1 + x^7 + x^{22} - x^{55};$$

`realroot(f3);`

$$f4 := -10 + 2x^7 + 10x^{22} - 20x^{55};$$

`realroot(f4)`

$$f1 := -x^{55} + x^{22} - x^7 - 1$$

$$\left[\left[-\frac{30953}{32768}, -\frac{61903}{65536} \right] \right]$$

$$f2 := -x^{55} + 10x^{22} - 2x^7 - 2$$

$$\left[\left[-\frac{118191}{131072}, -\frac{472761}{524288} \right], \left[\frac{61}{64}, \frac{977}{1024} \right], \left[\frac{139993}{131072}, \frac{279989}{262144} \right] \right]$$

$$f3 := -x^{55} + x^{22} + x^7 - 1$$

$$\left[[-1, -1], \left[\frac{125029}{131072}, \frac{250335}{262144} \right], [1, 1] \right]$$

$$f4 := -20x^{55} + 10x^{22} + 2x^7 - 10$$

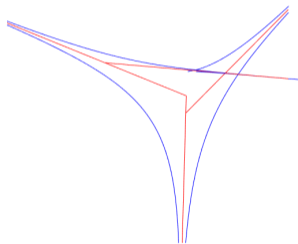
$$\left[\left[-\frac{1001}{1024}, -\frac{125}{128} \right] \right]$$

Figure: Using Maple software

Now we can see which region future polynomials lie in which will give us the number of real roots!

Canonical slice of $\text{Nabla}_A(\mathbb{R})$, plotted on log paper, for the family

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With many thanks...

- Thank you Dr. Rojas for the Matlab code!
- Thank you to Dr. Rojas and TA Joshua Goldstein for the guidance!
- Thank you to the NSF and Texas A & M for making this research experience possible!